## CMSC 330: Organization of Programming Languages

Lets, Tuples, Records

## Let Expressions

- Syntax
  - -let x = e1in e2
  - x is a bound variable
  - e1 is the binding expression
  - e2 is the body expression
- let expressions bind *local* variables
  - Different from let definitions, which are at the top-level

## Let Expressions

- Syntax
  - -let x = e1in e2

- Evaluation
  - $-e1 \Rightarrow v1$
  - $e2{v1/x}$

let z = 3+4 in 3\*z

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## Let Expressions

- Syntax
  - -let x = e1in e2
- Type checking
  - If e1: t1 and
  - If assuming x : t1 implies e2 : t
  - Then (let x = e1 in e2): t

#### Example

What is the type of let z = 3+4 in 3\*z?

- 3+4:int
- Assuming z:int, we have 3\*z:int
- So the type of let z = 3+4 in 3\*z is int

## Let Definitions vs. Let Expressions

At the top-level, we write

```
- let x = e;; (* no in e2 part *)
```

- This is called a let definition, not a let expression
  - Because it doesn't, itself, evaluate to anything
- Omitting in means "from now on":

```
# let pi = 3.14;;
(* pi is now bound in the rest of the top-level scope *)
```

## Let Expressions: Scope

• In let x = e1 in e2, var x is not visible outside of e2

```
float pi = 3.14;

pi * 3.0 * 3.0;
}
pi; /* pi unbound! */
```

## Examples – Scope of Let bindings

```
• x;; (* Unbound value x *)
• let x = 1 in x + 1;; (*2*)
• let x = x in x + 1; (* Unbound value x*)
• (let x = 1 in x + 1);; x;; (* Unbound value x*)
• let x = 4 in (let x = x + 1 in x);; (* 5 *)
```

## Nested Let Expressions

```
let res =
  (let area =
        (let pi = 3.14 in
        let r = 3.0 in
        pi *. r *. r) in
        area /. 2.0);;
```

#### Similar scoping possibilities C and Java

```
float res;
{ float area;
    { float pi = 3.14
        float r = 3.0;
        area = pi * r * r;
    }
    res = area / 2.0;
}
```

## Let Expressions in Functions

You can use let inside of functions for local vars

```
let area d =
  let pi = 3.14 in
  let r = d /. 2.0 in
  pi *. r *. r
```

## **Shadowing Names**

- Shadowing is rebinding a name in an inner scope to have a different meaning
  - May or may not be allowed by the language

```
let x = 10 in
    let z =
        let x = 20 in
        x*2 in
    x+z. (* 50 *)
```

## Shadowing, by the Semantics

- What if e2 is also a let for x?
  - Substitution will stop at the e2 of a shadowing x

```
Example
    let x = 3+4 in let x = 3*x in x+1
    - let x = 7 in let x = 3*x in x+1
    - let x = 3*7 in x+1
    - let x = 21 in x+1
    - 21+1
    - 22
    Will not be substituted, since it is shadowed by the inner let
```

#### Quiz 1: What does this evaluate to?

```
let x = 2 in
let y = x + x in
y * x
```

- A. 4
- B. 6
- C. 8
- D. Error

#### Quiz 1: What does this evaluate to?

```
let x = 2 in
let y = x + x in
y * x
```

- A. 4
- B. 6
- C. 8
- D. Error

#### Quiz 2: What does this evaluate to?

let 
$$x = 5$$
 in  $x = 3$ 

- A. 3
- B. 2
- C. true
- D. false

#### Quiz 2: What does this evaluate to?

- A. 3
- B. 2
- C. true
- D. false

This expression is checking whether **x** is equal to 3

#### Quiz 3: What does this evaluate to?

```
let y = 3 in
let x = y+2 in
let y = 6 in
x+y
```

- A. 8
- B. 11
- C. 13
- D. 14

### Quiz 3: What does this evaluate to?

```
let y = 3 in
let x = y+2 in
let y = 6 in
x+y
```

- A. 8
- B. 11
- C. 13
- D. 14

## **Tuples**

- Constructed using (e1, ..., en)
- Deconstructed using pattern matching
  - Patterns involve parens and commas, e.g., (p1, p2, ...)
- Tuples are similar to C structs
  - But without field labels
  - Allocated on the heap
- Tuples can be heterogenous
  - Unlike lists, which must be homogenous
  - (1, ["string1";"string2"]) is a valid tuple

## Tuple Types

- Tuple types use \* to separate components
  - Type joins types of its components
- Examples

```
- (1, 2) :
- (1, "string", 3.5) :
- (1, ["a"; "b"], 'c') :
- [(1,2)] :
- [(1, 2); (3, 4)] :
- [(1,2); (1,2,3)] :
```

## Tuple Types

- Tuple types use \* to separate components
  - Type joins types of its components
- Examples

```
- (1, 2):

- (1, "string", 3.5):

- (1, ["a"; "b"], 'c'):

- [(1,2)]:

- [(1,2)]:

- [(1,2); (3, 4)]:

- [(1,2); (1,2,3)]:

Because the first list element has type int * int, but the second has type int * int * int - list elements
```

must all be of the same type

## Pattern Matching Tuples

```
let plus3 t =
  match t with
  (x, y, z) -> x + y + z;;
plus3 : int*int*int -> int = <fun>
let plus3' (x, y, z) = x + y + z;;
plusThree' : int*int*int -> int = <fun>
```

## Tuples Are A Fixed Size

This OCaml definition

```
- let foo x = match x with
    (a, b) -> a + b
| (a, b, c) -> a + b + c
```

has a type error. Why?

Tuples of different size have different types

```
- (a, b) has type: 'a * 'b
- (a, b, c) has type: 'a * 'b * 'c
```

#### Quiz 4: What does this evaluate to?

```
let get a b = (a+b,0) in
get 1 2
```

- A. (3,0)
- B. (2,0)
- C. 3
- D. type error

#### Quiz 4: What does this evaluate to?

```
let get a b = (a+b,0) in
get 1 2
```

- A. (3,0)
- B. (2,0)
- C. 3
- D. type error

#### Quiz 5: What does this evaluate to?

```
let get (a,b) y = a+y in get (2,1) 1
```

- A. 3
- B. type error
- C. 2
- D. 1

#### Quiz 5: What does this evaluate to?

```
let get (a,b) y = a+y in get (2,1) 1
```

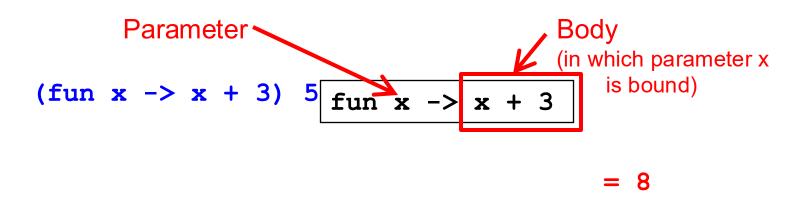
- A. 3
- B. type error
- C. 2
- D. 1

# CMSC 330 Organization of Programming Languages

## OCaml Higher Order Functions

## **Anonymous Functions**

Use fun to make a function with no name



## **Anonymous Functions**

- Syntax
  - fun *x1* ... *xn* -> *e*
- Evaluation
  - An anonymous function is an expression
  - In fact, it is a value.
- Type checking
  - (fun  $x1 \dots xn$  -> e):(t1 ->  $\dots$  -> tn -> u)
    when e: u under assumptions  $x1:t1, \dots, xn:tn$ .
    - (Same rule as let  $f \times 1 \dots \times n = e$ )

#### Quiz 1: What does this evaluate to?

```
let y = (fun x -> x+1) 2 in (fun z -> z-1) y
```

- A. Error
- B. 2
- C. 1
- D. 0

#### Quiz 1: What does this evaluate to?

```
let y = (fun x -> x+1) 2 in (fun z -> z-1) y
```

- A. Error
- B. 2
- C. 1
- D. 0

## Quiz 2: What is this expression's type?

$$(fun x y -> x) 2 3$$

- A. Type error
- B. int
- C. int -> int -> int
- D. 'a -> 'b -> 'a

## Quiz 2: What is this expression's type?

$$(fun x y -> x) 2 3$$

- A. Type error
- B. int
- C. int -> int -> int
- D. 'a -> 'b -> 'a

## **Functions and Binding**

 Functions are first-class, so you can bind them to other names as you like

```
let f x = x + 3;;
let g = f; # 8
g 5
```

## **Example Shorthands**

let for functions is a syntactic shorthand

```
let f x = body is semantically equivalent to
let f = fun x -> body
```

- let next x = x + 1
  - Short for let next =  $fun x \rightarrow x + 1$
- let plus x y = x + y
  - Short for let plus =  $fun x y \rightarrow x + y$

#### Quiz 3: What does this evaluate to?

```
let f = fun x -> 0 in
let g = f in
let h = fun y -> g (y+1) in
h 1
```

- **A**. 0
- B. 1
- C. 2
- D. Error

#### Quiz 3: What does this evaluate to?

```
let f = fun x -> 0 in
let g = f in
let h = fun y -> g (y+1)
h 1
```

- **A**. 0
- B. 1
- C. 2
- D. Error

#### **Nested Functions**

```
(* Filter the odd numbers from a list *)
let filter lst =
  let rec aux 1 =
       match 1 with
      |[] -> []|
      |h::t-> if h mod 2 <> 0 then h::aux t
        else aux t
      in
  aux 1st
filter [1;2;3;4;5;6] (* int list = [1; 3; 5] *)
```

## Passing Functions as Arguments

You can pass functions as arguments

```
let plus3 x = x + 3 (* int -> int *)
let twice f z = f (f z)
(* ('a->'a) -> 'a -> 'a *)

twice plus3 5 11
```

#### Records

- Records: identify elements by name
  - Elements of a tuple are identified by position
- Define a record type before defining record values

```
type date = { month: string; day: int; year: int }
```

Define a record value

```
# let today = { day=16; year=2017; month="f"^"eb" };;
today : date = { day=16; year=2017; month="feb" };;
```

## **Destructing Records**

```
type date = { month: string; day: int; year: int }
let today = { day=16; year=2017; month="feb" };;
```

Access by field name or pattern matching

```
today.month;; (* feb *)
let { year } = today in (* binds year to 2017 *)
let { month=_; day=d } = today in
...
```

## Quiz 6: What is the type of shift?

```
type point = {x:int; y:int}
let shift { x = px } = [px]::[]
```

```
A. point -> int listB. int -> int listC. point -> point listD. point -> int list list
```

## Quiz 6: What is the type of shift?

```
type point = {x:int; y:int}
let shift { x = px } = [px]::[]
```

```
A. point -> int listB. int -> int listC. point -> point listD. point -> int list list
```