



## **Problem 1: Basics**

[Total 4 pts]

There are some data structures in Rust which will not deallocate using the Reference Counting Garbage Collection Strategy	True	False F
There are some data structures in Rust which will deallocate using the Reference Counting Garbage Collection Strategy	T	F
Rust's Type System prevents Double Frees unless the <b>unsafe</b> keyword is used	T	F
Rust's Type System prevents Double Frees unless the <b>safe</b> keyword is used	T	F
It is theoretically possible to implement project 3 (NFA to DFA) in Lambda Calculus	T	F
It is theoretically possible to implement project 4 (MicroCaml) in Lambda Calculus	T	F
It is impossible to implement project 4 (MicroCaml) in Lambda Calculus	T	F
It is impossible to implement project 3 (NFA to DFA) in Lambda Calculus	T	F
$(\lambda x.y)((\lambda x.xx)(\lambda z.zz))$ has a beta normal form under eager evaluation	T	F
$(\lambda x.x)((\lambda y.yy)(\lambda z.zz))$ has a beta normal form under eager evaluation	T	F
$(\lambda y.y)((\lambda y.xy)(\lambda z.zz))$ has a beta normal form under eager evaluation	T	F
$(\lambda y.x)((\lambda x.xx)(\lambda y.xy))$ has a beta normal form under eager evaluation	T	F

## Problem 2: Lambda Calculus - Variables

[Total 2 pts]

Underline the <u>free variables</u> and circle the <u>bound variables</u>) in the expression below.

**Note**: Do not mark any of the lambda parameter variables.

### **Version A:**

$$\underline{a}(\lambda a. \lambda b. \underline{b} \lambda a. \underline{a})(\lambda c. \underline{d}) \underline{c}$$

# **Version B:**

$$\underline{z}(\lambda f.(\lambda b.\underline{a}\lambda a.\underline{a}))(\lambda c.\underline{c})\underline{c}$$

## **Version C:**

$$(\lambda a. \lambda b. b) (\lambda a. a) \underline{z} \lambda c. \underline{d} \underline{d}$$

#### **Version D:**

$$(\lambda x. \lambda b. \cancel{x}) (\lambda y. \underline{a} \cancel{y} \lambda c. \cancel{c}) \underline{d}$$

# Problem 3: Lambda Calculus - Alpha Equivalence

[Total 2 pts]

#### **Version A:**

Which lambda calculus expressions are alpha equivalent to  $(\lambda a. a)((\lambda b. c \lambda x. x) a b c)$ ? Circle all that apply.

- $(\lambda a. a)((\lambda a. c \lambda a. a) a b c)$   $(B)(c \lambda a. a) c$

- $(C)(\lambda c. a)((\lambda b. c \lambda c. c) a b c) \qquad (D)(\lambda f. f)((\lambda c. c \lambda g. g) a b c)$

#### **Version B:**

Which lambda calculus expressions are alpha equivalent to  $(\lambda b. a)((\lambda c. c \lambda b. b) \times y z)$ ? Circle all that apply.

- $(A)(\lambda a. a)((\lambda b. c \lambda a. a) a b c)$
- $(\lambda x. a)((\lambda g. g \lambda a. a) x y z)$
- $(\lambda d. a)((\lambda c. c \lambda b. b) \times y z) \qquad (D)(\lambda b. b)((c \lambda a. a) b c)$

### **Version C:**

Which lambda calculus expressions are alpha equivalent to  $(\lambda a.\ b)((\lambda b.\ c\ \lambda c.\ c)\ a\ b\ c)$ ? Circle all that apply.

- $(A)(\lambda x. b)((\lambda c. c \lambda d. d) a b c) \qquad (B)(\lambda x. a)((\lambda g. g \lambda a. a) x y z)$
- $(C)(\lambda a. b)((\lambda c. c \lambda c. c) a b c) \qquad (D)(\lambda a. b)((\lambda x. c \lambda c. c) a b c)$

#### **Version D:**

Which lambda calculus expressions are alpha equivalent to  $(\lambda b. b)((\lambda y. z \lambda d. d) x y z)$ ? Circle all that apply.

- $(\lambda c. c)((\lambda d. z \lambda d. d) \times y z)$
- $(\lambda b. b)((\lambda f. z \lambda f. f) \times y z)$
- $(C)(\lambda a. b)((\lambda c. c \lambda c. c) a b c) \qquad (D)(\lambda a. b)((\lambda x. c \lambda c. c) a b c)$

# Problem 4: Lambda Calculus - Reduction

[Total 4 pts]

Reduce the given lambda expression to beta normal form and show each step.

#### **Version A:**

Reduce  $(\lambda a.(\lambda b.(\lambda c.c c)b)a)d$ 

 $(\lambda b.(\lambda c.c c)b)d$  $(\lambda c.c c)d$ (d d)

### **Version C:**

Reduce  $(\lambda x.(\lambda y.(\lambda z.z z)y)x)a$  $(\lambda y.(\lambda z.zz)y)a$  $(\lambda z.zz)a$ (a a)

#### **Version B:**

Reduce  $(\lambda a.(\lambda b.(\lambda c.c f)b)a)x$ 

 $(\lambda b.(\lambda c.c f)b)x$  $(\lambda c.c f)x$ (x f)

#### **Version D:**

Reduce  $(\lambda x.(\lambda y.(\lambda z.z f)y)x)d$  $(\lambda y.(\lambda z.z f)y)d$  $(\lambda z.z f)d$ (df)

## Version A, C:

```
fn main(){
  {
    let a = String::from("hello");
    let b = f1(a);
    // Mark 1
    let c = f2(\&b);
    // Mark 2
  }
  // Mark 3
fn f1(s: String) -> String{
  println!("{}",s.len());
  // Mark 4
}
fn f2(s: &str)-> i32{
    s.len() as i32
Version B, D:
fn main(){
    let a = String::from("hello");
    // Mark 1
    let b = f1(a);
    // Mark 2
    let c = f2(\&b);
    // Mark 3
  }
}
fn f1(s: String) -> String{
  println!("{}",s.len());
  // Mark 4
fn f2(s: &str)-> i32{
    s.len() as i32
}
```

If there is no owner (because the value has been dropped) put "None". Assume that we are asking about ownership **during** execution.

Who is the owner of the value "hello" at Mark 1?

b

Who is the owner of the value "hello" at Mark 2?

b

Who is the owner of the value "hello" at Mark 3?

None

Who is the owner of the value "hello" at Mark 4?

S

If there is no owner (because the value has been dropped) put "None". Assume that we are asking about ownership **during** execution.

Who is the owner of the value "hello" at Mark 1?

a

Who is the owner of the value "hello" at Mark 2?

b

Who is the owner of the value "hello" at Mark 3?

b

Who is the owner of the value "hello" at Mark 4?

s/b