



## CMSC330 Fall 2024 Quiz 3 Solutions

### Problem 1: Context Free Grammars - Derivations

[Total 6 pts]

Consider the following Grammar:

$$\begin{aligned} S &\rightarrow FSH|q \\ F &\rightarrow Ff|f \\ H &\rightarrow hHh|\epsilon \end{aligned}$$

(a) Derive ffqhh

[4 pts]

$S \rightarrow FSH \rightarrow FfSH \rightarrow ffSH \rightarrow ffqH \rightarrow ffqhHh \rightarrow ffqhh$

(b) Is this an ambiguous grammar?

[2 pts]

A Yes       B No

$$\begin{aligned} S &\rightarrow ASB|c \\ A &\rightarrow aA|a \\ B &\rightarrow bbB|\epsilon \end{aligned}$$

(c) Derive aacbb

[4 pts]

$S \rightarrow ASB \rightarrow aASB \rightarrow aaSB \rightarrow aacB \rightarrow aacbbB \rightarrow aacbb$

(d) Is this an ambiguous grammar?

[2 pts]

A Yes       B No

$$\begin{aligned} S &\rightarrow CSD|f \\ C &\rightarrow cC|c \\ D &\rightarrow Ddd|\epsilon \end{aligned}$$

(e) Derive ccfdd

[4 pts]

$S \rightarrow CSD \rightarrow cCSD \rightarrow ccSD \rightarrow ccfD \rightarrow ccfDdd \rightarrow ccfdd$

(f) Is this an ambiguous grammar?

[2 pts]

A Yes       B No

$$\begin{aligned} S &\rightarrow NSP|s \\ N &\rightarrow Nn|n \\ P &\rightarrow pPp|\epsilon \end{aligned}$$

(g) Derive nnspp

[4 pts]

$S \rightarrow NSP \rightarrow NnSP \rightarrow nnSP \rightarrow nnsP \rightarrow nnsPP \rightarrow nnspp$

(h) Is this an ambiguous grammar?

[2 pts]

A Yes       B No

## Problem 2: Lexing Parsing and Evaluating

[Total 6 pts]

Given the following CFG, and assuming the **Ocaml** type system and semantics, at what stage of language processing would each expression fail? Mark '**Valid**' if the expression would be accepted by the grammar and evaluate successfully. Assume the only symbols allowed are those found in the grammar.

$$E \rightarrow M \text{ and } E|M \text{ or } E|M$$

$$M \rightarrow N + M | N - M | N$$

$$N \rightarrow 1 | 2 | 3 | 4 | \text{true} | \text{false} | (E)$$

	<b>Lexer</b>	<b>Parser</b>	<b>Evaluator</b>	<b>Valid</b>
1 + 2 - (true and false)	L	P	E	V
{2}	L	P	E	V
3 * 1 - 2	L	P	E	V
2 and 5	L	P	E	V
false	L	P	E	V
true and (false)	L	P	E	V
true - (false and 2)	L	P	E	V
(5)	L	P	E	V
3 / (1 - 2)	L	P	E	V
2 + 3 - 4 or true	L	P	E	V
true or true	L	P	E	V
(true and false)	L	P	E	V
1 and 2 - (true - false)	L	P	E	V
true * 1	L	P	E	V
1 - (true or false)	L	P	E	V
true	L	P	E	V
true and (false or 3)	L	P	E	V
1 + true and false	L	P	E	V
4 / 2	L	P	E	V
(2 - true) and false	L	P	E	V
2 and 1	L	P	E	V
0 - 1	L	P	E	V
true or (false - 2)	L	P	E	V

### Problem 3: Operational Semantics

[Total 4 pts]

Consider the following rules for two languages. OCaml will be the meta-language for both. Take note of the order of  $e_1$  and  $e_2$  that is bolded in Language B.

#### LANGUAGE A

$$\begin{array}{c} (\text{int rule}) \frac{}{A; n \Rightarrow n} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 * v_2}{A; \text{op1 } e_1 \ e_2 \Rightarrow v_3} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 + v_2}{A; \text{op2 } e_1 \ e_2 \Rightarrow v_3} \end{array}$$

#### LANGUAGE B

$$\begin{array}{c} (\text{int rule}) \frac{}{A; n \Rightarrow n} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 * v_2}{A; \mathbf{e}_2 \ e_1 \ \text{op3} \Rightarrow v_3} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 + v_2}{A; e_1 \ e_2 \ \text{op4} \Rightarrow v_3} \end{array}$$

Assume we derive meaning through operational semantics and read rules left-to-right. Give the Language B sentence that is semantically the same as the Language A sentence:

op1 3 op2 6 4

6 4 op4 3 op3

#### LANGUAGE A

$$\begin{array}{c} (\text{int rule}) \frac{}{A; n \Rightarrow n} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 * v_2}{A; \text{op1 } e_1 \ e_2 \Rightarrow v_3} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 + v_2}{A; \text{op2 } e_1 \ e_2 \Rightarrow v_3} \end{array}$$

#### LANGUAGE B

$$\begin{array}{c} (\text{int rule}) \frac{}{A; n \Rightarrow n} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 * v_2}{A; e_1 \ e_2 \ \text{op3} \Rightarrow v_3} \\[10pt] \dfrac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 + v_2}{A; e_2 \ e_1 \ \text{op4} \Rightarrow v_3} \end{array}$$

Assume we derive meaning through operational semantics and read rules left-to-right. Give the Language B sentence that is semantically the same as the Language A sentence:

op2 5 op1 8 1

8 1 op3 5 op4

## LANGUAGE A

$$\begin{array}{c}
 (\text{int rule}) \frac{}{A; n \Rightarrow n} \\
 \hline
 \begin{array}{c}
 A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 * v_2 \\
 \hline
 A; \text{op1 } e_1 \ e_2 \Rightarrow v_3
 \end{array} \\
 \hline
 \begin{array}{c}
 A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 + v_2 \\
 \hline
 A; \text{op2 } e_1 \ e_2 \Rightarrow v_3
 \end{array}
 \end{array}$$

Assume we derive meaning through operational semantics and read rules left-to-right. Give the Language B sentence that is semantically the same as the Language A sentence:

op1 op2 4 3 2

2 4 3 op4 op3

## LANGUAGE A

$$\begin{array}{c}
 (\text{int rule}) \frac{}{A; n \Rightarrow n} \\
 \hline
 \begin{array}{c}
 A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 * v_2 \\
 \hline
 A; \text{op1 } e_1 \ e_2 \Rightarrow v_3
 \end{array} \\
 \hline
 \begin{array}{c}
 A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 + v_2 \\
 \hline
 A; \text{op2 } e_1 \ e_2 \Rightarrow v_3
 \end{array}
 \end{array}$$

Assume we derive meaning through operational semantics and read rules left-to-right. Give the Language B sentence that is semantically the same as the Language A sentence:

op1 op2 8 6 7

6 8 op4 7 op3

## Problem 4: Type Checking

[Total 4 pts]

Consider the following type checking rules of OCaml:

$$\begin{array}{c}
 \frac{}{G \vdash n : \text{int}} \quad \frac{G(x) = t}{G \vdash x : t} \\[10pt]
 \frac{G \vdash e_1 : \text{int} \quad G \vdash e_2 : \text{int} \quad (+) : \text{int} \rightarrow \text{int} \rightarrow \text{int}}{G \vdash e_1 + e_2 : \text{int}} \\[10pt]
 \frac{G \vdash e_1 : \text{int} \quad G \vdash e_2 : \text{int} \quad (*) : \text{int} \rightarrow \text{int} \rightarrow \text{int}}{G \vdash e_1 * e_2 : \text{int}} \\[10pt]
 \frac{G \vdash e_1 : \text{int} \quad G \vdash e_2 : \text{int} \quad (-) : \text{int} \rightarrow \text{int} \rightarrow \text{int}}{G \vdash e_1 - e_2 : \text{int}} \\[10pt]
 \frac{G \vdash e_1 : t_1 \quad G, x : t_1 \vdash e_2 : t_2}{G \vdash \text{let } x = e_1 \text{ in } e_2 : t_2}
 \end{array}$$

Write a type checking proof for the following expression

$$\begin{array}{c}
 \text{let } x = 5 \text{ in } x + 7 \\[10pt]
 \frac{\frac{\frac{G, x : \text{int}(x) = \text{int}}{G, x : \text{int} \vdash x : \text{int}} \quad \frac{G, x : \text{int} \vdash 7 : \text{int}}{(+) : \text{int} \rightarrow \text{int} \rightarrow \text{int}}}{G, x : \text{int} \vdash x + 7 : \text{int}}}{G \vdash \text{let } x = 5 \text{ in } x + 7}
 \end{array}$$

Write a type checking proof for the following expression

$$\begin{array}{c}
 \text{let } x = 6 \text{ in } 1 + x \\[10pt]
 \frac{\frac{\frac{G, x : \text{int}(x) = \text{int}}{G, x : \text{int} \vdash 1 : \text{int}} \quad \frac{G, x : \text{int} \vdash x : \text{int}}{(+) : \text{int} \rightarrow \text{int} \rightarrow \text{int}}}{G, x : \text{int} \vdash 1 + x : \text{int}}}{G \vdash \text{let } x = 6 \text{ in } 1 + x}
 \end{array}$$

Write a type checking proof for the following expression

$$\begin{array}{c}
 \text{let } x = 2 \text{ in } x * 9 \\[10pt]
 \frac{\frac{\frac{G, x : \text{int}(x) = \text{int}}{G, x : \text{int} \vdash x : \text{int}} \quad \frac{G, x : \text{int} \vdash 9 : \text{int}}{(*) : \text{int} \rightarrow \text{int} \rightarrow \text{int}}}{G, x : \text{int} \vdash x * 9 : \text{int}}}{G \vdash \text{let } x = 2 \text{ in } x * 9}
 \end{array}$$

Write a type checking proof for the following expression

$$\begin{array}{c}
 \text{let } x = 42 \text{ in } 0 - x \\[10pt]
 \frac{\frac{\frac{G, x : \text{int}(x) = \text{int}}{G, x : \text{int} \vdash 0 : \text{int}} \quad \frac{G, x : \text{int} \vdash x : \text{int}}{(-) : \text{int} \rightarrow \text{int} \rightarrow \text{int}}}{G, x : \text{int} \vdash 0 - x : \text{int}}}{G \vdash \text{let } x = 42 \text{ in } 0 - x}
 \end{array}$$