

CMSC330 - Organization of Programming Languages Fall 2024 - Final

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Name: _____

UID: _____

I pledge on my honor that I have not given or received any unauthorized assistance on this assignment/examination

Signature: _____

Ground Rules

- Please write legibly. **If we cannot read your answer you will not receive credit.**
- You may use anything on the accompanying reference sheet anywhere on this exam
- Please remove the reference sheet from the exam
- The back of the reference sheet has some scratch space on it. If you use it, you must turn in your scratch work
- You may not leave the room or hand in your exam within the last 10 minutes of the exam
- If anything is unclear, ask a proctor. If you are still confused, write down your assumptions in the margin

Question	Points
P1.	10
P2.	10
P3.	15
P4.	6
P5.	4
P6.	6
P7.	6
P8.	6
P9.	10
P10.	3
P11.	8
P12.	16
Total	100

Problem 1: Concepts

[Total 10 pts]

	True	False
Some Buffer Overflow vulnerabilities can be prevented by having a type-safe type system	<input checked="" type="radio"/>	<input type="radio"/>
In Rust, null pointer exceptions will not occur	<input checked="" type="radio"/>	<input type="radio"/>
If you are at some state B in an FSM, the history of your path determines where you go next	<input type="radio"/>	<input checked="" type="radio"/>
Context Free Grammars can recognize strings with balanced parenthesis	<input checked="" type="radio"/>	<input type="radio"/>
The LL(k) parser we used in project 4 runs in little- $o(n)$ time where n is the number of tokens.	<input checked="" type="radio"/>	<input type="radio"/>
Stop and Copy Garbage collection will not clean up cyclic data structures	<input type="radio"/>	<input checked="" type="radio"/>
The rules of references that Rust uses helps prevents double frees	<input type="radio"/>	<input checked="" type="radio"/>
Both Ocaml and Rust are statically typed	<input checked="" type="radio"/>	<input type="radio"/>
The best case runtime of the tokenize function from project 4 is polynomial.	<input checked="" type="radio"/>	<input type="radio"/>
In Rust, a reference's lifetime is part of its type	<input checked="" type="radio"/>	<input type="radio"/>
OCaml is Turing complete which means it can solve more problems than Rust	<input type="radio"/>	<input checked="" type="radio"/>
If a language is well typed, it must also be well-defined	<input type="radio"/>	<input checked="" type="radio"/>
In Rust, a reference's lifetime is not part of its type	<input type="radio"/>	<input checked="" type="radio"/>
If a language is well-typed, it is possible for it to not be well-defined	<input checked="" type="radio"/>	<input type="radio"/>
Both Ocaml and Rust are dynamically typed	<input type="radio"/>	<input checked="" type="radio"/>
Rust is Turing complete which means it can solve more problems than Ocaml	<input type="radio"/>	<input checked="" type="radio"/>
Regular Expressions can recognize strings of arbitrary length with balanced parenthesis	<input type="radio"/>	<input checked="" type="radio"/>

Problem 2: Regex

[Total 10 pts]

If you run `ping google.com -c 2` on the command line, you get an echo response from your destination sending 2 network packets. It can be useful to see if you have an active internet connection or if a website is down. An example ping command can return:

```
PING google.com (192.168.255.255) 56 bytes of data
64 bytes from 192.168.255.255: icmp_seq=1 ttl=60 time=8.57 ms
64 bytes from 192.168.255.255: icmp_seq=2 ttl=60 time=2.61 ms
--- google.com ping statistics ---
2 packet transmitted, 2 received, 0% packet loss, time 1003ms
```

Write a regex that describes each part of this echoed response:

(a) IP Address

[2 pts]

IP Addresses appear multiple times in the following lines. 192.168.255.255 is an IP address. Valid IP addresses will look like 'xxx.xxx.xxx.xxx' where 'xxx' is from 0-255 (inclusive). **Write the regex for the 'xxx' part (the range of 0-255, inclusive).** This regex will be used in the later parts as **IP**.

```
([0-1][0-9][0-9]|2[0-4][0-9]|25[0-5])
```

(b) Destination Summary - **This is the first line in the example**

[2 pts]

Example: `PING google.com (192.168.255.255) 56 bytes of data`

It summarizes where you are pinging and how many bytes the sent data will be. The domain name will be **at least one lowercase letter** followed by any number of dots (.) and lowercase letters with no consecutive dots (it may end in a dot). Only 56 or 64 bytes of data will be sent.

```
[a-z](\.[a-z]+)*\.(56|64)
```

(c) Received Data

[2 pts]

The second and third lines in the example are instances of this. Received data will receive 32 or 64 bytes from an IP address. `icmp_seq` is the sequence number of the packet (≥ 0), `ttl` will also be ≥ 0 . Time will be any number ≥ 0 with 2 digits after the decimal. You will not need to check if the sequence is in order, or if the IP address is consistent across responses.

```
(32|64) bytes from IP: icmp_seq=[0-9]+ ttl=[0-9]+ time=[0-9]+\.[0-9]{2} ms
```

(d) Statistics

[3 pts]

The last line in the example. The statistics line will have how many packets were transmitted (≥ 0), how many were received (≥ 0), the percent packet loss (0 – 100 inclusive), and the time (≥ 0). Time in this response does not include decimals and milliseconds. You don't need to check if the math is right.

```
[0-9]+ packet transmitted, [0-9]+ received ([0-9][0-9]?100)% packet loss, time [0-9]+ms
```

If you run `ping google.com -c 2` on the command line, you get an echo response from your destination sending 2 network packets. It can be useful to see if you have an active internet connection or if a website is down. An example ping command can return:

```
PING google.com (567.451.701.701) 56 bytes of data
64 bytes from 567.451.701.701: ttl=60 icmp_seq=1 time=8.57 ms
64 bytes from 567.451.701.701: ttl=60 icmp_seq=2 time=2.61 ms
--- google.com ping statistics ---
2 packet transmitted, 2 received, 0% packet loss, time 1003ms
```

Write a regex that describes each part of this echoed response:

(e) IP Address

[2 pts]

IP Addresses appear multiple times in the following lines. 567.168.255.990 is an IP address. For our purposes, valid IP addresses will look like 'xxx.xxx.xxx.xxx' where 'xxx' is from 400-742 (inclusive). Write the regex for the IP address here and then **in each following location where it should appear write: IP**

```
(([4-6][0-9][0-9]|7[0-3][0-9]|74[0-2])\.([4-6][0-9][0-9]|7[0-3][0-9]|74[0-2])){3}
```

(f) Destination Summary

[2 pts]

This is the first line in the example. It summarizes where you are pinging and how many bytes the sent data will be. The domain name will be **at least one lowercase letter** followed by any number of dots (.) and lowercase letters with no repeating dots (it may end in a dot). Only 56 or 64 bytes of data will be sent.

```
PING [a-z](\.[a-z]+) * \.? IP(64|56) bytes of data
```

(g) Received Data

[2 pts]

The second and third lines in the example are instances of this. Received data will receive 32 or 64 bytes from an IP address. ttl is (≥ 0), icmp_seq will also be ≥ 0 . Time will be any number ≥ 0 with 2 digits after the decimal. You will not need to check if the sequence is in order, or if the IP address is consistent across responses.

```
(32|64) bytes from IP: ttl=[0-9]+ icmp_seq=[0-9]+ time=[0-9]+\.[0-9]{2}ms
```

(h) Statistics

[3 pts]

The last line in the example. The statistics line will have how many packets were transmitted (≥ 0), how many were received (≥ 0), the percent packet loss (0 – 100 inclusive), and the time (≥ 0). Time in this response does not include decimals and milliseconds. You don't need to check if the math is right.

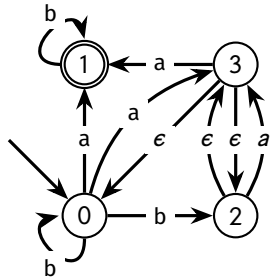
```
[0-9]+ packet transmitted, [0-9]+ received ([0-9][0-9]?|100)% packet loss, time [0-9]+ms
```

Problem 3: FSM

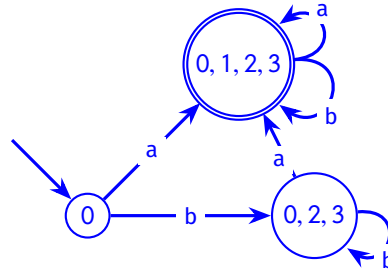
[Total 15 pts]

(a) Convert the below NFA to a DFA.
Draw a **box** around your final answer.

[10 pts]



Scratch Space:



(b) Write a **CFG** that describes strings accepted by the NFA above.

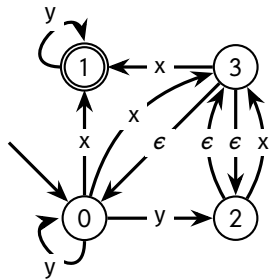
[5 pts]

$$S \rightarrow bS \mid aT$$

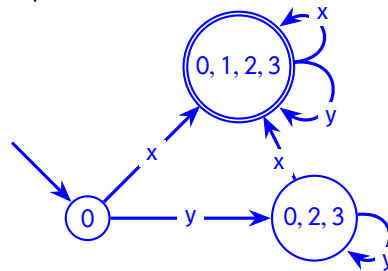
$$T \rightarrow aT \mid bT \mid \epsilon$$

(c) Convert the below NFA to a DFA.
Draw a **box** around your final answer.

[10 pts]



Scratch Space:



(d) Write a **CFG** that describes strings accepted by the NFA above.

[5 pts]

$$S \rightarrow yS \mid xT$$

$$T \rightarrow xT \mid yT \mid \epsilon$$

Problem 4: OCaml Typing

[Total 5 pts]

Give the type of the variable 'r'. If there is a type error, put "ERROR"

```
let e a b =  
  fold_left (fun x y -> a x y) true (1.2 :: b)
```

```
(bool -> float -> bool) -> float list -> bool
```

```
let e =  
  fun y ->  
  fun z ->  
  map y (2 :: z)
```

```
(int -> 'b) -> int list -> 'b list
```

```
let foo a = fun b -> map a (map b [1;2;3])
```

```
('a -> 'b) -> (int -> 'a) -> 'b list
```

```
let foo a b c = if a (b c) then  
  c  
  else  
  c
```

```
('b -> bool) -> ('a -> 'b) -> 'a -> 'a
```

Problem 5: Evaluation

[Total 6 pts]

Evaluate the following OCaml expressions. If there is a compilation error, put "ERROR"

```
let foo a = fun b -> map a (map b [1;2;3]) in  
foo (fun x -> -x) (fun y -> y * 4)
```

```
[-4;-8;-12]
```

```
let foo a b c = if a (b c) then  
  c  
  else  
  c  
in foo (fun a -> 0) (fun b -> b) true
```

```
ERROR
```

```
let e a b =  
  fold_left (fun x y -> a x y) true (1.2 :: b)  
in  
e (fun x -> x > 2.5) [3.14]
```

```
ERROR
```

```
let e =  
  fun y ->  
  fun z ->  
  map y (2 :: z)  
in  
e (fun x -> x + 1) [4; 3]
```

```
[3;5;4]
```

Problem 6: Property Based Testing

[Total 10 pts]

Consider an attempted (buggy!) implementation of the `double_up_and_dance` function from project 6. `double_up_and_dance` *should* create a vector that contains duplicates of each item in the input slice, with the second item in the slice added to the fifth if the duplicated vector is longer than 4 elements. It then returns the created vector.

```
pub fn double_up_and_dance(slice: &[i32]) -> Vec<i32> {
    let mut ret = Vec::with_capacity(slice.len() * 2);
    for &elem in slice {
        ret.push(elem);
        ret.push(elem);
    }
    if let Some(val) = ret.get_mut(5) {
        *val += slice[0];
    }
    ret
}
```

Consider the property p :

The 5th element of the `double_up_and_dance` vector will be greater than the 2nd element of the input slice (if they exist).

Is p a valid property? Yes No

Suppose we wanted to write this test. We would encode the property as the following:

```
fn test_prop(v in prop::collection::vec(usize,1..10)){
    //will generate random usize vectors of lengths 1 through 10
    if v.len() < 3{
        assert!(true)
    }else{
        assert!(double_up_and_dance(&v).get(5).unwrap() > v.get(2).unwrap())
    }
}
```

Is `test_prop` a correct encoding of the property p ? Yes No

If we test this property on the provided implementation of `double_up_and_dance`, will it ever assert false?

Yes

No

Problem 7: Interpreters

[Total 6 pts]

Given the following CFG, at what stage of language processing would each expression **fail**?

Mark **'Valid'** if the expression would be accepted by the grammar and evaluate properly. Assume the only symbols allowed are those found in the grammar. Choose only one choice for each expression.

Note: For expressions that result in an infinite loop, consider them to fail at the evaluator step.

$$E \rightarrow LX.E \mid EE \mid (E) \mid X$$

$$X \rightarrow c$$

For all $c \in X$, c is a lowercase English character

	Lexer	Parser	Evaluator	Valid
LZ.a	<input checked="" type="radio"/> L	<input type="radio"/> P	<input type="radio"/> E	<input type="radio"/> V
Lx.(Ly.x y)	<input type="radio"/> L	<input type="radio"/> P	<input type="radio"/> E	<input checked="" type="radio"/> V
x ((Ld.e y) e)	<input type="radio"/> L	<input checked="" type="radio"/> P	<input type="radio"/> E	<input type="radio"/> V
(La.(Ly.(e e) Ly.(y y)))	<input type="radio"/> L	<input type="radio"/> P	<input type="radio"/> E	<input checked="" type="radio"/> V
(Lt.(Lx.(t y x)).z)	<input type="radio"/> L	<input checked="" type="radio"/> P	<input type="radio"/> E	<input type="radio"/> V
(Lx. x x) (Lx. x x)	<input type="radio"/> L	<input type="radio"/> P	<input checked="" type="radio"/> E	<input type="radio"/> V
Lx.a	<input type="radio"/> L	<input type="radio"/> P	<input type="radio"/> E	<input checked="" type="radio"/> V
(Lx. x x) (La. a a)	<input type="radio"/> L	<input type="radio"/> P	<input checked="" type="radio"/> E	<input type="radio"/> V
Lx.y.(Ly.x y)	<input type="radio"/> L	<input checked="" type="radio"/> P	<input type="radio"/> E	<input type="radio"/> V
(La.(Ly1.(e e) Ly2.(y2 y2)))	<input checked="" type="radio"/> L	<input type="radio"/> P	<input type="radio"/> E	<input type="radio"/> V
(Lt.(Lx.(t (y x))) z)	<input type="radio"/> L	<input type="radio"/> P	<input type="radio"/> E	<input checked="" type="radio"/> V

Problem 8: Type Checking

[Total 16 pts]

Consider the following Typing Rules for Ocaml:

$$\begin{array}{c}
 \overline{G \vdash \text{true} : \text{bool}} \\
 \\
 \overline{G \vdash x : G(x)} \\
 \\
 \frac{G \vdash e_1 : t_1 \quad G, x : t_1 \vdash e_2 : t_2}{G \vdash \text{let } x = e_1 \text{ in } e_2 : t_2} \\
 \\
 \overline{G \vdash \text{false} : \text{bool}} \quad \overline{G \vdash n : \text{int}} \\
 \\
 \frac{G \vdash e_1 : \text{int} \quad G \vdash e_2 : \text{int} \quad += (\text{int}, \text{int}, \text{int})}{G \vdash e_1 + e_2 : \text{int}} \\
 \\
 \frac{G \vdash e_1 : \text{bool} \quad G \vdash e_2 : t \quad G \vdash e_3 : t}{G; \text{if } e_1 \text{ then } e_2 \text{ else } e_3 : t}
 \end{array}$$

Write a type-checking proof for the expression

let x = true in if x then 4 else 5 + 7

$$\frac{\overline{G \vdash \text{true} : \text{bool}} \quad \frac{\overline{G, x : \text{bool} \vdash x : G(x)} \quad \overline{G, x : \text{bool} \vdash 4 : \text{int}} \quad \frac{\overline{G, x : \text{bool} \vdash 5 : \text{int}} \quad \overline{G, x : \text{bool} \vdash 7 : \text{int}} \quad += (\text{int}, \text{int}, \text{int})}{\overline{G, x : \text{bool} \vdash 5 + 7 : \text{int}}}}{\overline{G, x : \text{bool} \vdash \text{if } x \text{ then } 4 \text{ else } 5 + 7 : \text{int}}} \\
 \overline{G \vdash \text{let } x = \text{true in if } x \text{ then } 4 \text{ else } 5 + 7 : \text{int}}$$

Write a type-checking proof for the expression

if let x = false in if x then 9 + 1 else 2

$$\frac{\overline{G \vdash \text{false} : \text{bool}} \quad \frac{\overline{G, x : \text{bool} \vdash x : G(x)} \quad \overline{G, x : \text{bool} \vdash 9 : \text{int}} \quad \overline{G, x : \text{bool} \vdash 1 : \text{int}} \quad += (\text{int}, \text{int}, \text{int})}{\overline{G, x : \text{bool} \vdash 9 + 1 : \text{int}}} \quad \overline{G, x : \text{bool} \vdash 2 : \text{int}}}{\overline{G, x : \text{bool} \vdash \text{if } x \text{ then } 9 + 1 \text{ else } 2 : \text{int}}} \\
 \overline{G \vdash \text{if let } x = \text{false in if } x \text{ then } 9 + 1 \text{ else } 2 : \text{int}}$$

Problem 9: Lambda Calculus

[Total 10 pts]

(a) Reduce

[6 pts]

Reduce the following lambda expression. Show every step.

$$(\lambda c. c d (\lambda a. a))(\lambda b. (\lambda c. (\lambda a. d b)))$$

$$\begin{aligned} &(\lambda c. c d (\lambda a. a))(\lambda b. (\lambda c. (\lambda a. d b))) \\ &((\lambda b. (\lambda c. (\lambda a. d b))) d (\lambda a. a)) \\ &((\lambda c. (\lambda a. d d)) (\lambda a. a)) \\ &(\lambda a. d d) \end{aligned}$$

$$(\lambda x. (\lambda c. cx)) (cc) (\lambda b. (\lambda a. a b))$$

$$\begin{aligned} &(\lambda x. (\lambda c. cx)) (cc) (\lambda b. (\lambda a. ab)) \\ &(\lambda a. (\lambda x. xa)) (cc) (\lambda b. (\lambda a. ab)) \\ &(\lambda x. x(cc)) (\lambda b. (\lambda a. ab)) \\ &(\lambda b. (\lambda a. ab)) (cc) \\ &(\lambda a. a(cc)) \end{aligned}$$

(b) Free Variables:

[2 pts]

Circle the free variables in the expression below:

$$(\lambda a. (\lambda c. a \boxed{b})) ((\lambda b. (\lambda a. b)) (\boxed{a} \boxed{a} (\lambda a. \boxed{c})))$$
$$((\lambda a. (\lambda c. a)) (\lambda b. (\lambda a. b \boxed{c})) (\boxed{a} \boxed{a} (\lambda a. \boxed{c})))$$

(c) Alpha Equivalence:

[2 pts]

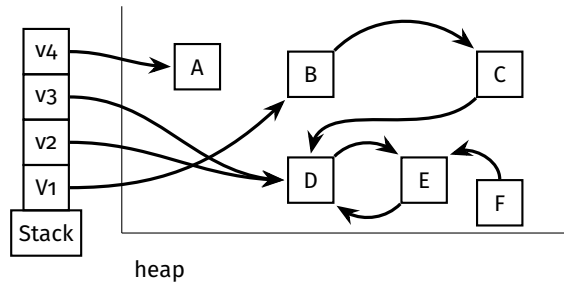
Which of the following are alpha equivalent to the expression $(\lambda b. (\lambda a. b)(ca))$? Select all that apply.

- (A) $(\lambda z. (\lambda y. z)(xy))$
- (B) $(\lambda a. (\lambda a. a)(ca))$
- (C) $(\lambda a. (\lambda b. a)(ca))$
- (D) $(\lambda b. (\lambda c. b)(ca))$

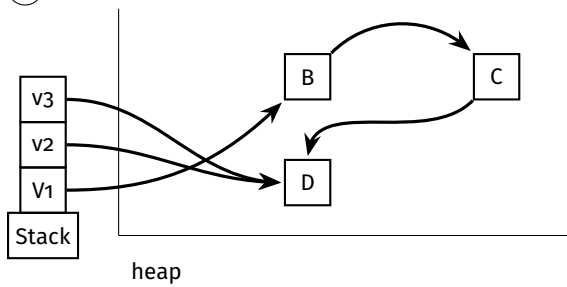
Problem 10: Garbage Collection

[Total 3 pts]

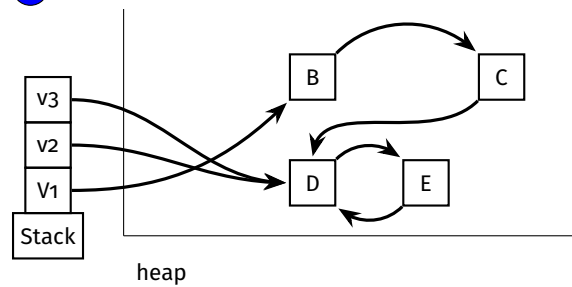
Suppose you have the following memory diagram. Select the correct representation of the stack and heap after the mark and sweep garbage collection technique is run after popping v_4 off the stack.



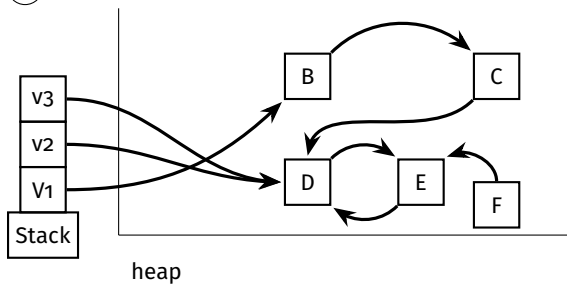
(A)



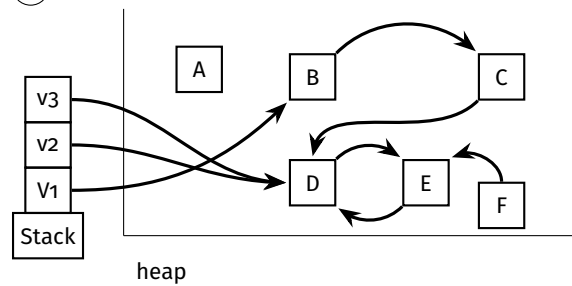
(B)



(C)



(D)



(E) None of the above

Problem 11: Ownership and Lifetimes

[Total 8 pts]

Does the code compile? Yes No

```
fn main(){
    let mut x = 4;
    let y = &mut x;
    println!("{x},{y}");
}
```

If **no**, explain why not in one sentence:

println cannot borrow immut while y is borrow mut

Does the code compile? Yes No

```
fn main(){
    let x = String::from("bye");
    let y = String::from("farewell");

    let mut z = &x;
    z = y;
    println!("{x},{z}")
}
```

If **no**, explain why not in one sentence:

Does the code compile? Yes No

```
fn main(){
    let x = String::from("hello");
    let mut y = x;
    y.push_str(" world!");
    println!("{y}");
}
```

If **no**, explain why not in one sentence:

```
struct Thing{
    v1:usize,
    v2:String
}
```

Does the code compile? Yes No

```
fn main(){
    let mut x = Thing{v1:4,v2:String::from("hi")};
    let y = &mut x;
    y.v2.push_str(" bye");
    let z = y.v2;
    println!("{},{z}",x.v1)
}
```

If **no**, explain why not in one sentence:

y.v2 tries to move ownership of a borrowed value

Problem 12: Coding

[Total 18 pts]

(a) OCaml

[? pts]

Given an `int list list`, return true if the first column (the first element in each list) is in descending order, and false otherwise. Assume the matrix is guaranteed to have equal height and width ($m \times m$ matrix) and non-empty ($m > 0$). If items are equal, that is not descending (i.e. a column of `[1;1;0]` is not descending).

```
ex.                ex
[[1;2;3;];         [[7;8;9];
 [4;5;6]; => false  [4;5;6]; => true
 [7;8;9]]          [1;6;3]]
```

```
let rec decending mtx =
let f::t = map (fun x::_ -> x) mtx in
let r,_ = fold (fun (a,b) x -> (x,(x < a) && b)) (f,true) t in
r
```

```
let rec accending mtx =
let f::t = map (fun x::_ -> x) mtx in
let r,_ = fold (fun (a,b) x -> (x,(x > a) && b)) (f,true) t in
r
```

(b) Rust

[? pts]

Given an `Vec<Vec<u32>>`, return true if the first column (the first element in each list) is in descending order, and false otherwise. Assume the matrix is guaranteed to have equal height and width ($m \times m$ matrix) and non-empty ($m > 0$). If items are equal, that is not descending (i.e. a column of `[1;1;0]` is not descending).

```
ex.                ex
vec![vec![1,2,3],   vec![vec![7,8,9],
  vec![4,5,6], => false    vec![4,5,6], => true
  vec![7,8,9]]         vec![1,6,3]]
```

```
fn decending(v:Vec<Vec<u32>>)->bool{
  let mut r = v.get(0).unwrap().get(0).unwrap();
  for i in &v[1..]{
    if i.get(0).unwrap() < r{
      r = i.get(0).unwrap()
    } else{
      return false;
    }
  }
};
true
}
```

```
fn ascending(v:Vec<Vec<u32>>)->bool{
  let mut r = v.get(0).unwrap().get(0).unwrap();
  for i in &v[1..]{
    if i.get(0).unwrap() > r{
      r = i.get(0).unwrap()
    } else{
      return false;
    }
  }
};
true
}
```

Cheat Sheet

OCaml

```
(* Map and Fold *)
(* ('a -> 'b) -> 'a list -> 'b list *)
let rec map f l = match l with
  [] -> []
  | x::xs -> (f x)::(map f xs)

(* ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a *)
let rec fold_left f a l = match l with
  [] -> a
  | x::xs -> fold_left f (f a x) xs

(* ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b *)
let rec fold_right f l a = match l with
  [] -> a
  | x::xs -> f x (fold_right f xs a)
```

```
(* OCaml Function Types *)
:: -: 'a -> 'a list -> 'a list
```

```
@ -: 'a list -> 'a list -> 'a list
```

```
+, -, *, / -: int -> int -> int
+., -., *., /. -: float -> float -> float
```

```
&&, || -: bool -> bool -> bool
not -: bool -> bool
```

```
^ -: string -> string -> string
```

```
=>, >, =, <, <= :- 'a -> 'a -> bool
```

```
(* Regex in OCaml *)
```

```
Re.Posix.re: string -> regex
```

```
Re.compile: regex -> compiled_regex
```

```
Re.exec: compiled_regex -> string -> group
```

```
Re.exectp: compiled_regex -> string -> bool
```

```
Re.exec_opt: compiled_regex -> string -> group option
```

```
Re.matches: compiled_regex -> string -> string list
```

```
Re.Group.get: group -> int -> string
```

```
Re.Group.get_opt: group -> int -> string option
```

Structure of Regex

```
R → ∅
   | σ
   | ε
   | RR
   | R|R
   | R*
```

Regex

*	zero or more repetitions of the preceding character or group
+	one or more repetitions of the preceding character or group
?	zero or one repetitions of the preceding character or group
.	any character
$r_1 r_2$	r_1 or r_2 (eg. $a b$ means 'a' or 'b')
[abc]	match any character in abc
[$\sim r_1$]	anything except r_1 (eg. [$\sim abc$] is anything but an 'a', 'b', or 'c')
[r_1-r_2]	range specification (eg. [a-z] means any letter in the ASCII range of a-z)
{n}	exactly n repetitions of the preceding character or group
{n,}	at least n repetitions of the preceding character or group
{m,n}	at least m and at most n repetitions of the preceding character or group
^	start of string
\$	end of string
(r_1)	capture the pattern r_1 and store it somewhere (match group in Python)
\d	any digit, same as [0-9]
\s	any space character like \n, \t, \r, \f, or space

NFA to DFA Algorithm (Subset Construction Algorithm)

NFA (input): $(\Sigma, Q, q_0, F_n, \delta)$, DFA (output): $(\Sigma, R, r_0, F_d, \delta_n)$

```
 $R \leftarrow \{\}$ 
 $r_0 \leftarrow \epsilon - \text{closure}(\sigma, q_0)$ 
while  $\exists$  an unmarked state  $r \in R$  do
  mark  $r$ 
  for all  $a \in \Sigma$  do
     $E \leftarrow \text{move}(\sigma, r, a)$ 
     $e \leftarrow \epsilon - \text{closure}(\sigma, E)$ 
    if  $e \notin R$  then
       $R \leftarrow R \cup \{e\}$ 
    end if
     $\sigma_n \leftarrow \sigma_n \cup \{r, a, e\}$ 
  end for
end while
 $F_d \leftarrow \{r \mid \exists s \in r \text{ with } s \in F_n\}$ 
```

Rust

```
// Vectors
let vec = Vec::new(); // makes a new vector
let vec1 = vec![1,2,3]

vec.push(ele); // Pushes the element 'ele'
               // to end of the vector 'vec'

vec.get(x); // returns the xth index of the
            // vec in an option, with Some(ele)
            // Some(ele) if it exists
            // and None if it doesn't exist

// Strings
let string = String::from("Hello");

string.push_str(&str); // appends the str
                     // to string

vec.to_iter(); // returns an iterator for vec

string.chars() // returns an iterator of chars
              // over the a string
```

```
iter.rev(); // reverses an iterators direction

iter.next(); // returns an Option of the next
             // item in the iterator.

struct Building{ // example of struct
  name:String,
  floors:i32,
  locationx:f32,
  locatiny:f32,
}

enum Option<T>{ Some(T); None } //enum Option type
option.unwrap(); // returns the item in an Option or
                // panics if None

if let x = option {...} // assigns x to the element
                       // in the option if it is Some
```